

# Graph (1)

- Consist of a set of nodes and a set of edges between those nodes

- Most important model for applied mathematics

- Incidence matrix  $\mathbf{A}$  ( $m \times n$ )

( $m$  edges and  $n$  nodes)

edge  $i = \text{row } i$ , node  $j \rightarrow k : -1$  in column  $j$ ,  $+1$  in column  $k$

$N(\mathbf{A})$  contains all constant vectors:  $\mathbf{x} = (c, c, \dots, c)$

$\dim N(\mathbf{A}) = 1$ ,  $\dim C(\mathbf{A}) = \dim C(\mathbf{A}^T) = n - 1$ ,  $\dim N(\mathbf{A}^T) = m - (n - 1)$

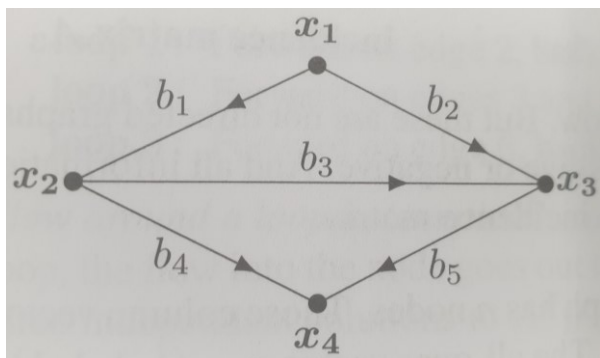
row space contains all constant vectors  $\mathbf{x}$  with  $x_1 + x_2 + \dots + x_n = 0$  ( $\mathbf{x} \perp \mathbf{1}$ )

continuous	discrete
function	vector
derivative	difference
integral	sum
calculus	linear algebra

bases  $\left\{ \begin{array}{l} N(\mathbf{A}): \text{ constant vector } \mathbf{1} \\ C(\mathbf{A}^T): (n-1) \text{ rows of } \mathbf{A} \text{ that produce a tree in the graph (a tree has no loop)} \\ C(\mathbf{A}): \text{ any } (n-1) \text{ columns of } \mathbf{A} \\ N(\mathbf{A}^T): \text{ flows around the } (m-n+1) \text{ small loops in the graph} \end{array} \right.$

# Graph (2)

- Example:  $m=5$ ,  $n=4$



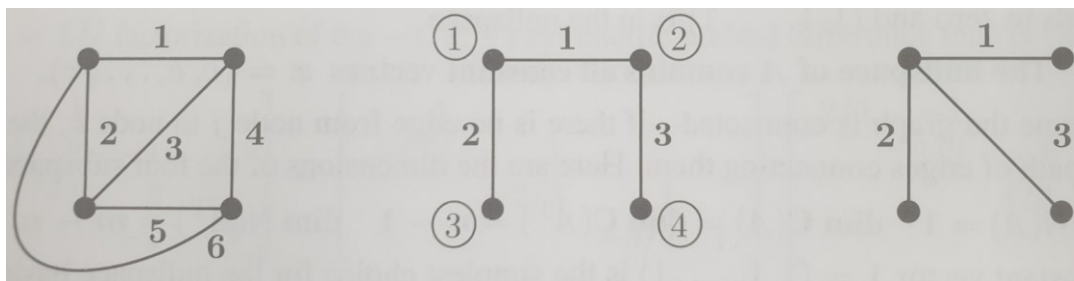
$$\begin{array}{r}
 \text{edges} \\
 \mathbf{A} = \begin{bmatrix} -1 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & -1 & 0 & 1 \\ 0 & 0 & -1 & 1 \end{bmatrix} \begin{array}{l} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{array} \\
 \text{nodes} \quad 1 \quad 2 \quad 3 \quad 4
 \end{array}$$

- Graph Laplacian matrix  $\mathbf{L} = \mathbf{A}^T \mathbf{A}$ 
  - Symmetric, positive semidefinite

$$\mathbf{L} = \mathbf{A}^T \mathbf{A} = \begin{bmatrix} 2 & -1 & -1 & 0 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ 0 & -1 & -1 & 2 \end{bmatrix} = \underbrace{\begin{bmatrix} 2 & & & \\ & 3 & & \\ & & 3 & \\ & & & 2 \end{bmatrix}}_{\substack{\text{degree matrix} \\ \text{count edges into node}}} - \underbrace{\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}}_{\substack{\text{adjacency matrix} \\ b_{jk}=1: \text{edge from } j \text{ to } k}} = \mathbf{D} - \mathbf{B}$$

# Graph (2)

- Complete graph: every pair of nodes is connected by an edge,  $D=(n-1)I$ ,  $B=\text{all-ones minus } I$
- Tree: there are no loops in the connected graph



$$\mathbf{A}_1 = \begin{bmatrix} -1 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & -1 & 0 & 1 \\ 0 & 0 & -1 & 1 \\ -1 & 0 & 0 & 1 \end{bmatrix}, \quad \mathbf{A}_2 = \begin{bmatrix} -1 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ 0 & -1 & 0 & 1 \end{bmatrix}, \quad \mathbf{A}_3 = \begin{bmatrix} -1 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ -1 & 0 & 0 & 1 \end{bmatrix}$$

complete graph:  $m = \frac{1}{2}n(n-1)$

→ any graph:  $(n-1) \leq m \leq \frac{1}{2}n(n-1)$

# Kirchhoff's Current Law: $A^T y = f$

- KCL = balance of currents (forces, money)
  - Flow into each node equals flow out from that node
  - Key to solving  $A^T y = 0$  is to look at the small loops in the graph
  - $(m-n+1)$  independent solutions
  - $(\text{number of nodes}) - (\text{number of edges}) + (\text{number of loops}) = 1$

$$\mathbf{A}_1 = \begin{bmatrix} -1 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & -1 & 0 & 1 \\ 0 & 0 & -1 & 1 \\ -1 & 0 & 0 & 1 \end{bmatrix} \rightarrow \underbrace{\mathbf{A}_1 \mathbf{x} = \mathbf{0}}_{\substack{\mathbf{x}=(1,1,1,1) \\ \text{not interesting!}}}, \mathbf{A}_1^T \mathbf{y} = \begin{bmatrix} -1 & -1 & 0 & 0 & 0 & -1 \\ 1 & 0 & -1 & -1 & 0 & 0 \\ 0 & 1 & 1 & 0 & -1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ y_4 \\ y_5 \\ y_6 \end{bmatrix} = \mathbf{0} \rightarrow \mathbf{y}_1 = \begin{bmatrix} -1 \\ 1 \\ -1 \\ 0 \\ 0 \\ 0 \end{bmatrix}, \mathbf{y}_2 = \begin{bmatrix} 0 \\ 0 \\ 1 \\ -1 \\ 1 \\ 0 \end{bmatrix}, \mathbf{y}_3 = \begin{bmatrix} 0 \\ -1 \\ 0 \\ 0 \\ -1 \\ 1 \end{bmatrix}$$

$$\text{outer loop} = \mathbf{y}_1 + \mathbf{y}_2 + \mathbf{y}_3$$

$$\mathbf{A}_2^T \mathbf{y} = 0 \rightarrow \mathbf{y} = 0$$

# A<sup>T</sup>CA Framework in Applied Mathematics

- Graphs are perfect examples for three equations in engineering, science, economics
  - Describe a system in steady state equilibrium
  - Balance laws: conservation of charge, balance of force, zero net income in economics, conservation of mass and energy, continuity of every kind

$\left\{ \begin{array}{l} \text{voltages } \mathbf{x} = (x_1, x_2, x_3, x_4) \text{ at the four nodes} \\ \text{currents } \mathbf{y} = (y_1, y_2, y_3, y_4, y_5, y_6) \text{ at along the six edges} \end{array} \right.$

$\left\{ \begin{array}{lll} \text{Voltage differences across edges} & \mathbf{e} = \mathbf{Ax} & e_1 = (\text{voltage at end node 2}) - (\text{voltage at end node 1}) \\ \text{Ohm's law on each edge} & \mathbf{y} = \mathbf{Ce} & \text{current } y_1 = c_1 \text{ times } e_1 = (\text{conductance})(\text{voltage}) \\ \text{Kirchhoff's Law with current sources} & \mathbf{f} = \mathbf{A}^T \mathbf{y} & \text{current sources } \mathbf{f} \text{ into nodes balance the internal currents } \mathbf{y} \end{array} \right.$

$$\rightarrow \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{x} = \mathbf{f} \rightarrow \mathbf{K} \mathbf{x} = \mathbf{f}$$

$\mathbf{K}$ : symmetric, positive **semi**definite  $\xrightarrow[n-1=3 \text{ unknown voltages}]{\text{boundary condition } x_4=0}$  **reduced**  $\mathbf{K}$ : symmetric, invertible, positive definite  $(3 \times 6)(6 \times 6)(6 \times 3)$

$$\text{energy: } \mathbf{x}^T (\mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{x}) = (\mathbf{A} \mathbf{x})^T \mathbf{C} (\mathbf{A} \mathbf{x}) > 0 \text{ if } \mathbf{x} \neq \mathbf{0}$$

# $A^T C A$ Framework in Applied Mathematics

- Linear regression: least squares applied to  $Ax=b$

$$\left\{ \begin{array}{l} \mathbf{A}^T \mathbf{A} \hat{\mathbf{x}} = \mathbf{A}^T \mathbf{b} : \text{Normal equation for the vector } \hat{\mathbf{x}} \text{ that best fits the data } \mathbf{b} \\ \mathbf{A}^T \mathbf{C} \mathbf{A} \hat{\mathbf{x}} = \mathbf{A}^T \mathbf{C} \mathbf{b} : \text{Least squares weighted by the inverse covariance matrix } \mathbf{C} = \mathbf{V}^{-1} \\ \min \|\mathbf{b} - \mathbf{A} \mathbf{x}\|_{\mathbf{C}}^2 : \text{Minimum squared error } (\mathbf{b} - \mathbf{A} \mathbf{x})^T \mathbf{C} (\mathbf{b} - \mathbf{A} \mathbf{x}) \end{array} \right.$$

- Graph Laplacian Matrix

$\mathbf{K} = \mathbf{A}^T \mathbf{C} \mathbf{A}$  : weighted graph Laplacian,  $\mathbf{G} = \mathbf{A}^T \mathbf{A}$  : standard Laplacian ( $\mathbf{C} = \mathbf{I}$ )

$\mathbf{A}^T \mathbf{A} = (\text{diagonal}) + (\text{off-diagonal}) = (\text{degree matrix}) - (\text{adjacency matrix}) = \mathbf{D} - \mathbf{B}$

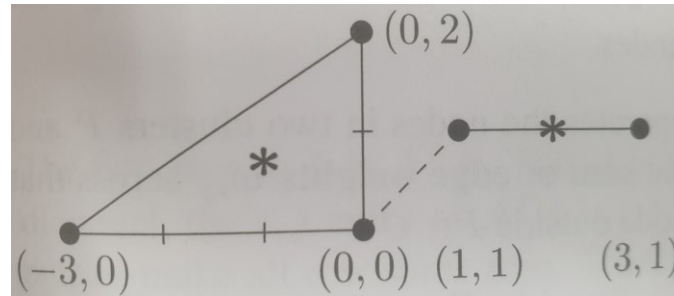
$\left\{ \begin{array}{l} \text{every row and column of } \mathbf{G} \text{ and } \mathbf{K} \text{ adds to zero because } \mathbf{x} = (1, \dots, 1) \text{ has } \mathbf{A} \mathbf{x} = \mathbf{0} \\ \mathbf{G} = \mathbf{A}^T \mathbf{A} \text{ is symmetric because edges go both ways (undirected graph)} \\ \text{The diagonal entry } (\mathbf{A}^T \mathbf{A})_{ii} \text{ counts the edges meeting at node } i: \text{the degree} \\ \text{The off-diagonal entry is } (\mathbf{A}^T \mathbf{A})_{ij} = -1 \text{ when an edge connects node } i \text{ and } j \\ \mathbf{G} \text{ and } \mathbf{K} \text{ are positive semidefinite but not positive definite (because } \mathbf{A} \mathbf{x} = \mathbf{0} \text{ in } \mathbf{1}) \end{array} \right.$

# Clustering

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- How to understand a graph with many nodes?
  - Separate nodes into two or more clusters
  - Human Genome project: cluster genes that show highly correlated
- Break a graph in two pieces: clusters of nodes
  - Each cluster should contain roughly half of the nodes
  - The number of edges between clusters should be relatively small
- Examples
  - For load balancing in high computing, assign equal work to two processors
  - For social networks, identify two distinct groups
  - Segment an image
  - Reorder rows and columns of a matrix to make off-diagonal blocks sparse

# Example with Two Clusters



$n = 5$  nodes,  $k = 2$  clusters

centroid  $*$  :  $\mathbf{c}_1 = (-1, 2/3)$ ,  $\mathbf{c}_2 = (2, 1) \leftarrow$  minimize the sum of squared distances  $\|\mathbf{c} - \mathbf{a}_j\|^2$

Approximate an  $m \times n$  matrix of  $\mathbf{A}$  by  $\mathbf{CR} = \underbrace{(m \times k)}_{\text{low rank}} \underbrace{(k \times n)}_{\substack{\text{only } k \text{ columns} \\ \text{centroids of clusters}} \quad \substack{\text{single 1 and} \\ \text{ } k-1 \text{ zeros}}}$

$R_{ij} = 1$  (or 0) if centroid  $i$  is closest (or not) to the point  $\mathbf{x}_j$

$$\mathbf{A} = \begin{bmatrix} 0 & 1 & 3 & 0 & -3 \\ 0 & 1 & 1 & 2 & 0 \end{bmatrix} \approx \begin{bmatrix} -1 & 2 & 2 & -1 & -1 \\ 2/3 & 1 & 1 & 2/3 & 2/3 \end{bmatrix}$$

$$\mathbf{A} \approx \mathbf{CR} = \begin{bmatrix} -1 & 2 \\ 2/3 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix}$$



# Four Methods for Clustering

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- Spectral clustering (Fiedler vector)
  - using the graph Laplacian or the modularity matrix
- Minimum cut
- Weighted k-means

# Four Methods for Clustering

I.  $\mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{z} = \lambda \mathbf{D} \mathbf{z} \rightarrow \mathbf{z}$ : **Fiedler vector**  $\begin{cases} \lambda_1 = 0 \rightarrow \mathbf{z}_1 = (1, \dots, 1) \\ \lambda_2 \rightarrow \mathbf{z}_2: + / - \text{ components indicate two clusters of nodes} \end{cases}$

II.  $\mathbf{A}^T \mathbf{C} \mathbf{A} \rightarrow \mathbf{M} = \mathbf{B} - \frac{1}{2m} \mathbf{d} \mathbf{d}^T$  where  $\mathbf{d}$ : degrees of the  $n$  nodes (number of edges adjacent to the nodes)  
modularity matrix

choose **eigenvector** that comes with the **largest eigenvalue of M**

III. Find the **minimum normalized cut** that separates the nodes in two clusters P and Q

weight across cut:  $links(P) = \sum w_{ij}$  for  $i$  in  $P$  and  $j$  not in  $P$

size of cluster:  $size(P) = \sum w_{ij}$  for  $i$  in  $P$

normalized cut weight:  $N_{cut}(P, Q) = \frac{links(P)}{size(P)} + \frac{links(Q)}{size(Q)} \xrightarrow{k\text{-cut}} N_{cut}(P_1, \dots, P_k) = \sum_{i=1}^k \frac{links(P_i)}{size(P_i)}$

minimize  $N_{cut}(P, Q)$ : good partition of the graph  $\rightarrow$  application: segmentation of images

IV. nodes in the graph:  $\mathbf{a}_1, \dots, \mathbf{a}_n$ , clusters  $P$  and  $Q$  have centers  $\mathbf{c}_P \left( = \frac{\sum \mathbf{a}_i}{|P|} \right)$  and  $\mathbf{c}_Q$

**minimize the total squared distance** from nodes to those centroids:  $E = \sum_{i \in P} \|\mathbf{a}_i - \mathbf{c}_P\|^2 + \sum_{i \in Q} \|\mathbf{a}_i - \mathbf{c}_Q\|^2$

# Spectral Clustering (1)

$$\mathbf{A}^T \mathbf{C} \mathbf{A} \xrightarrow[\text{the Laplacian}]{\mathbf{D} \text{ normalizes}} \mathbf{L} = \mathbf{D}^{-1/2} \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{D}^{-1/2} = \mathbf{I} - \mathbf{N} \quad \text{where } n_{ij} = \frac{w_{ij}}{\sqrt{d_i d_j}} \text{ (normalized weights)}$$

triangular graph:  $n = 3$  nodes,  $m = 3$  edges,  $c_1, c_2, c_3 = w_{12}, w_{13}, w_{23}$

$$\left. \begin{array}{l} \mathbf{A}^T \mathbf{C} \mathbf{A} = \mathbf{D} - \mathbf{W} \\ \left[ \begin{array}{ccc} w_{12} + w_{13} & -w_{12} & -w_{13} \\ -w_{21} & w_{21} + w_{23} & -w_{23} \\ -w_{31} & -w_{32} & w_{31} + w_{32} \end{array} \right] \end{array} \right\} \rightarrow \left\{ \begin{array}{l} \mathbf{L} = \mathbf{D}^{-1/2} \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{D}^{-1/2} \\ \left[ \begin{array}{ccc} 1 & -n_{12} & -n_{13} \\ -n_{21} & 1 & -n_{23} \\ -n_{31} & -n_{32} & 1 \end{array} \right] \end{array} \right.$$

$\mathbf{L} = \mathbf{I} - \mathbf{N}$  is like a correlation matrix in statistics

1.  $\mathbf{L}$  is symmetric positive semidefinite: orthogonal eigenvectors, all eigenvalues  $\lambda \geq 0$
2. The eigenvectors for  $\lambda = 0$  is  $\mathbf{u} = (\sqrt{d_1}, \dots, \sqrt{d_n})$ . Then  $\mathbf{L}\mathbf{u} = \mathbf{D}^{-1/2} \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{1} = 0$ .
3. The second eigenvector  $\mathbf{v}$  of  $\mathbf{L}$  minimizes the Rayleigh quotient on a subspace.

$$\left( \begin{array}{l} \lambda_2 = \text{smallest nonzero eigenvalue of } \mathbf{L} \rightarrow \min_{\substack{\text{subject to} \\ \mathbf{x}^T \mathbf{u} = 0}} \frac{\mathbf{x}^T \mathbf{L} \mathbf{x}}{\mathbf{x}^T \mathbf{x}} = \frac{\mathbf{v}^T \mathbf{L} \mathbf{v}}{\mathbf{v}^T \mathbf{v}} = \lambda_2 \text{ at } \mathbf{x} = \mathbf{v} \\ \text{upper bound for } \lambda_2, \text{ for any } \mathbf{x} \text{ orthogonal to the first eigenvector } \mathbf{u} = \mathbf{D}^{-1/2} \mathbf{1} \end{array} \right)$$

# Spectral Clustering normalized vs. unnormalized

$$\mathbf{L}\mathbf{v} = \mathbf{D}^{-1/2} \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{D}^{-1/2} \mathbf{v} = \lambda \mathbf{v} \xrightarrow[\text{normalized Fiedler vector}]{\mathbf{z} = \mathbf{D}^{-1/2} \mathbf{v}}$$

$$\left\{ \begin{array}{l} \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{z} = \lambda \mathbf{D} \mathbf{z} \text{ with } \mathbf{1}^T \mathbf{D} \mathbf{z} = 0 \\ \text{generalized eigenvalue problem} \\ \text{eigenvector for } \lambda = 0 \text{ is } \mathbf{1} \\ \text{the next eigenvector } \mathbf{z} \text{ is } \mathbf{D}\text{-orthogonal to } \mathbf{1} \\ \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{z} = \lambda_2 \mathbf{D} \mathbf{z} \end{array} \right.$$

$$\min_{\substack{\mathbf{x}^T \mathbf{L} \mathbf{x} \\ \text{subject to} \\ \mathbf{x}^T \mathbf{u} = 0}} \frac{\mathbf{x}^T \mathbf{L} \mathbf{x}}{\mathbf{x}^T \mathbf{x}} \xrightarrow{\mathbf{x} = \mathbf{D}^{1/2} \mathbf{y}} \min_{\substack{\mathbf{y}^T \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{y} \\ \text{subject to} \\ \mathbf{1}^T \mathbf{D} \mathbf{y} = 0}} \frac{\mathbf{y}^T \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{y}}{\mathbf{y}^T \mathbf{D} \mathbf{y}} = \frac{\sum \sum w_{ij} (y_i - y_j)^2}{\sum d_i y_i^2} = \lambda_2 \text{ at } \mathbf{y} = \mathbf{z}$$

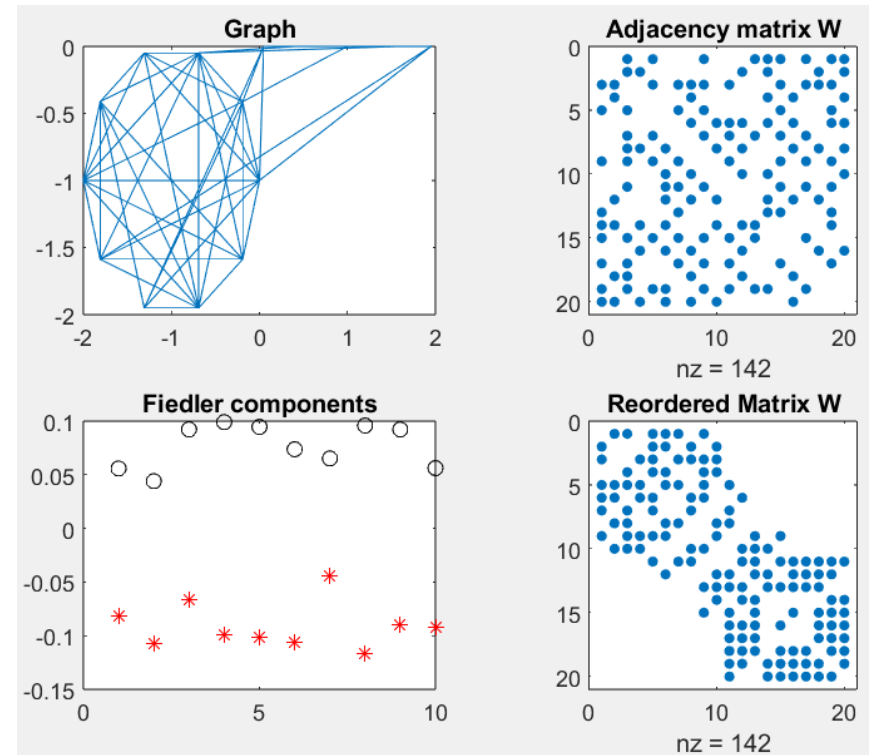
$\mathbf{A}\mathbf{y}$ : incidence matrix  $\mathbf{A}$  gives the differences  $(y_i - y_j)$

- Example: 20-node graph has two 10-node clusters P and Q (to find from  $\mathbf{z}$ )
  - Create edges within P and Q with probability 0.7. Edges between nodes in P and Q have smaller probability 0.1. All edges have weights  $w_{ij}=1$ . ( $\mathbf{C}=\mathbf{I}$ )

# Code: MATLAB

```
N=10; W=zeros(2*N,2*N); % Generate 2N nodes in two clusters
rand('state',100) % rand repeats to give the same graph
for i=1:2*N-1
for j=i+1:2*N
p=0.7-0.6*mod(j-i,2); % p=0.1 when j-i is odd, 0.7 else
W(i,j)=rand<p; % Insert edges with probability p
end % The weights are wi,j=1 (or 0)
end % So far W is strictly upper triangular
W=W+W'; D=diag(sum(W)); % Adjacency matrix W, degrees in D
G=D-W; [V,E]=eig(G,D); % Eigenvalues of Gx=(lambda)Dx in E
[a,b]=sort(diag(E)); z=V(:,b(2)); % Fiedler eigenvector z for (lambda)2
plot(sort(z), '-'); % Show +- groups of Fiedler components
```

```
theta=[1:N]*2*pi/N; x=zeros(2*N,1); y=x; % Angles to plot graph
x(1:2:2*N-1)=cos(theta)-1; x(2:2:2*N)=cos(theta)+1;
y(1:2:2*N-1)=sin(theta)-1; x(2:2:2*N)=sin(theta)+1;
print theta,x,y
subplot(2,2,1), gplot(W,[x,y]), title('Graph')
subplot(2,2,2), spy(W), title('Adjacency matrix W')
subplot(2,2,3), plot(z(1:2:2*N-1),'ko'), hold on
plot(z(2:2:2*N),'r*'), hold off, title('Fiedler components')
[c,d]=sort(z); subplot(2,2,4), spy(W(d,d)), title('Reordered Matrix W')
```



# Minimum Cut

(edge) weight across cut:  $links(P) = \sum w_{ij}$  for  $i$  in  $P$  and  $j$  not in  $P$

size of cluster:  $size(P) = \sum w_{ij}$  for  $i$  in  $P$

normalized cut weight:  $Ncut(P, Q) = \frac{links(P)}{size(P)} + \frac{links(Q)}{size(Q)}$

normalized  $K$ -cut:  $Ncut(P_1, \dots, P_k) = \sum_{i=1}^k \frac{links(P_i)}{size(P_i)}$

[cuts connected to eigenvectors]

perfect indicator of a cut: vector  $\mathbf{y}$  with all components equal to  $p$  or  $-q$  (two values only)  $\rightarrow$  node  $i$  goes  $\begin{cases} \text{in } P \text{ if } y_i = p \\ \text{in } Q \text{ if } y_i = -q \end{cases}$

$\mathbf{1}^T \mathbf{D} \mathbf{y}$  will multiply one group of  $d_i$  by  $p$  and the other group by  $-q$ .  
 The first  $d_i$  add to  $size(P) = \text{sum of } d_i \text{ (} i \text{ in } P)$ .  
 The second group of  $d_i$  add to  $size(Q)$

$$\left. \begin{array}{l} \mathbf{1}^T \mathbf{D} \mathbf{y} \text{ will multiply one group of } d_i \text{ by } p \text{ and the other group by } -q. \\ \text{The first } d_i \text{ add to } size(P) = \text{sum of } d_i \text{ (} i \text{ in } P). \\ \text{The second group of } d_i \text{ add to } size(Q) \end{array} \right\} \rightarrow \mathbf{1}^T \mathbf{D} \mathbf{y} = 0 \rightarrow psize(P) = qsize(Q)$$

$$\frac{\mathbf{y}^T \mathbf{A}^T \mathbf{C} \mathbf{A} \mathbf{y}}{\mathbf{y}^T \mathbf{D} \mathbf{y}} = \frac{\sum \sum w_{ij} (y_i - y_j)^2}{\sum d_i y_i^2} = \frac{(p+q)^2 links(P, Q)}{p^2 size(P) + q^2 size(Q)} = \frac{(p+q) links(P, Q)}{psize(P)} = \frac{links(P, Q)}{size(P)} + \frac{links(P, Q)}{size(Q)} = Ncut(P, Q)$$

# Clustering by k-means

$n$  points  $\mathbf{a}_1, \dots, \mathbf{a}_n$  in  $d$ -dimensional space  $\rightarrow$  partition those points into  $k$  clusters

clusters  $P_1, \dots, P_k$  have centroids  $\mathbf{c}_1, \dots, \mathbf{c}_k$

$\mathbf{c}_j = \frac{\text{sum of } \mathbf{a}'\text{s}}{\text{number of } \mathbf{a}'\text{s}} \rightarrow \text{minimize } \sum \|\mathbf{c} - \mathbf{a}\|^2$  for all  $\mathbf{a}$ 's in cluster  $P_j$

clustering: to find the partition  $P_1, \dots, P_k$  with minimum total distance  $D$  to centroids:

minimize  $D = \sum_{j=1}^k D_j = \sum_{j=1}^k \|\mathbf{c}_j - \mathbf{a}_i\|^2$  for  $\mathbf{a}_i$  in cluster  $P_j$

step 1: find the **centroids**  $\mathbf{c}_j$  of the (old) clustering  $P_1, \dots, P_k$ .

step 2: find the **(new) clustering** that puts  $\mathbf{a}$  in  $P_j$  if  $\mathbf{c}_j$  is the closest centroid.

# Clustering by k-means: Weights and Kernel Method

weighted distance:  $d(\mathbf{x}, \mathbf{a}_i) = w_i \|\mathbf{x} - \mathbf{a}_i\|^2$ ,  $\mathbf{c}_j = \frac{\sum w_i \mathbf{a}_i}{\sum w_i}$  ( $\mathbf{a}_i$  in  $P_j$ )

distances to centroids only require dot product  $\mathbf{a}_i \cdot \mathbf{a}_j$ : (each  $i$  in  $P_j$ )  $\|\mathbf{c}_j - \mathbf{a}_i\|^2 = \mathbf{c}_j \cdot \mathbf{c}_j - 2\mathbf{c}_j \cdot \mathbf{a}_i + \mathbf{a}_i \cdot \mathbf{a}_i$

Kernel method: weighted kernel matrix  $\mathbf{K}$  has entries  $\mathbf{a}_i \cdot \mathbf{a}_l$

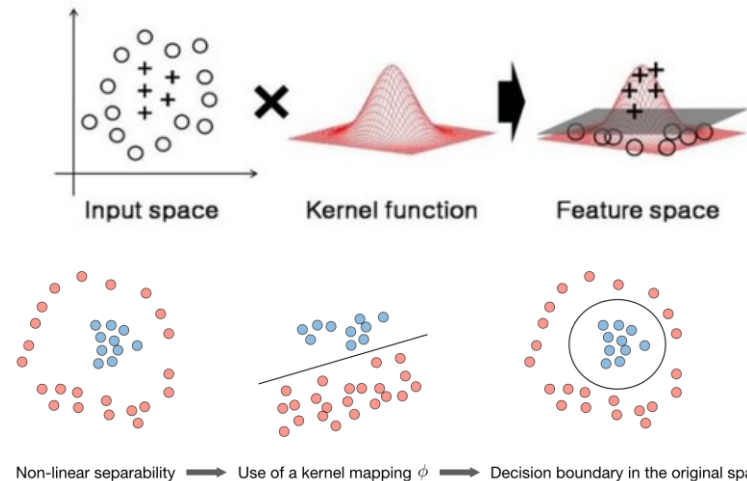
nodes are point  $\mathbf{x}_i$  in input space  $\rightarrow \mathbf{a}_i = \phi(\mathbf{x}_i)$  points in a high-dimensional **feature space**

$$\left(\text{sum over nodes in } P_j\right) \sum \|\mathbf{c}_j - \mathbf{a}_i\|^2 = \frac{\sum w_i w_l \mathbf{K}_{il}}{(\sum w_i)^2} - 2 \frac{\sum w_i \mathbf{K}_{il}}{\sum w_i} + \sum \mathbf{K}_{ii}$$

(vision) polynomial  $\mathbf{K}_{il} = (\mathbf{x}_i \cdot \mathbf{x}_l + c)^d$

(statistics) Gaussian  $\mathbf{K}_{il} = \exp\left(-\frac{\|\mathbf{x}_i - \mathbf{x}_l\|^2}{2\sigma^2}\right)$

(neural networks) Sigmoid  $\mathbf{K}_{il} = \tanh(c\mathbf{x}_i \cdot \mathbf{x}_l + \theta)$



for large data sets,  $k$ -means and  $\text{eig}(\mathbf{A}^T \mathbf{C} \mathbf{A}, \mathbf{D})$  will be expensive  $\rightarrow$   $\begin{cases} \text{random sampling} \\ \text{multilevel clustering} \end{cases}$



# Applications of Clustering

- Learning theory, training sets, neural networks, Hidden Markov Models
- Classification, regression, pattern recognition, Support Vector Machines
- Statistical learning, maximum likelihood, Bayesian statistics, spatial statistics, kriging, time series, ARMA models, stationary processes
- Social networks, organization theory
- Data mining, document indexing, image retrieval, kernel-based learning, Nystrom method, low rank approximation
- Bioinformatics, microarray data, systems biology
- Cheminformatics, drug design, decision trees
- Information theory, vector quantization, rate distortion theory, Bregman divergences
- Image segmentation, computer vision, texture, min cut
- Predictive control, feedback samples, robotics, adaptive control, Riccati equations